

Aspektorientierte Programmierung

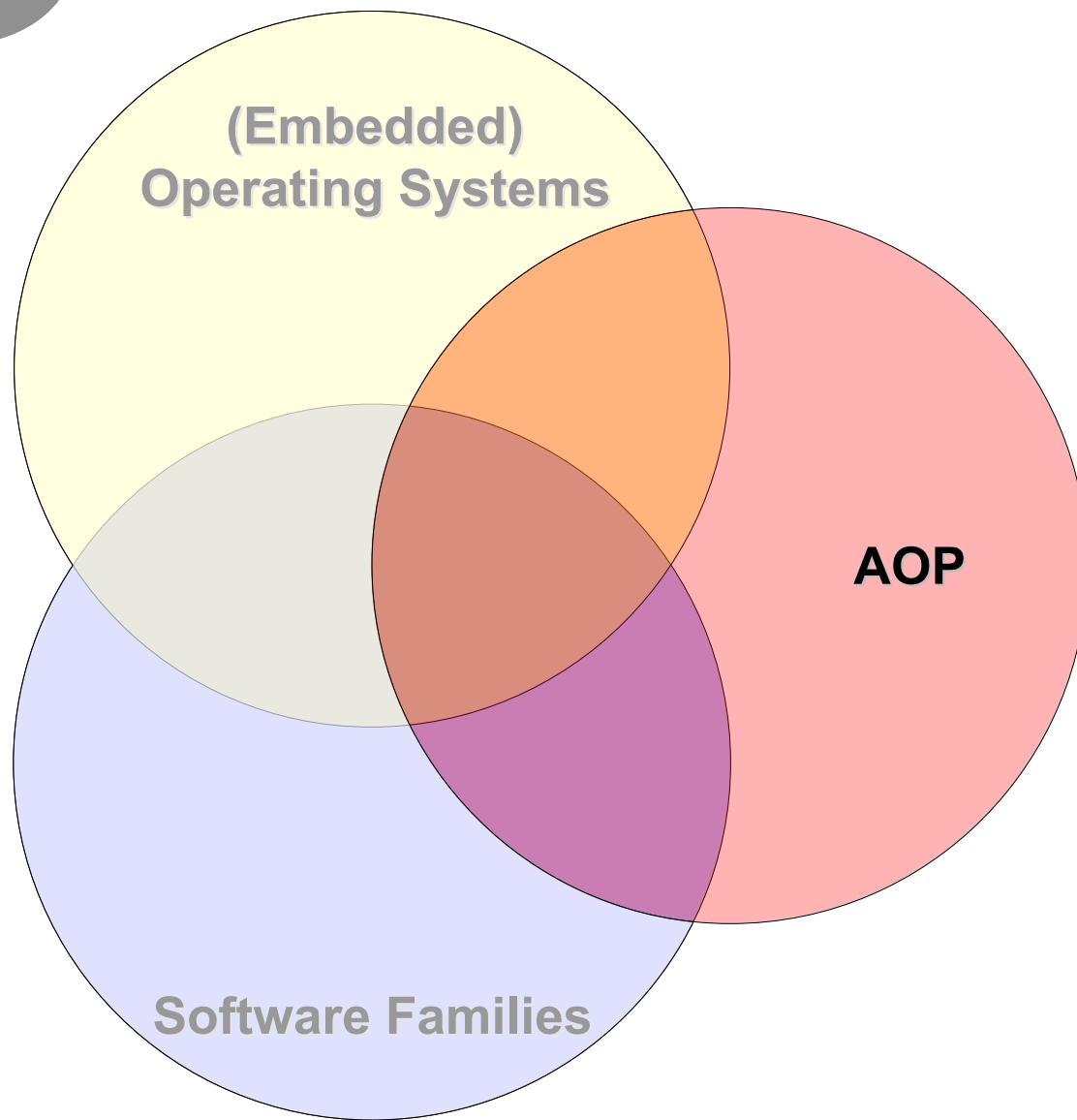
Eine Einführung am Beispiel von AspectC++

Daniel Lohmann

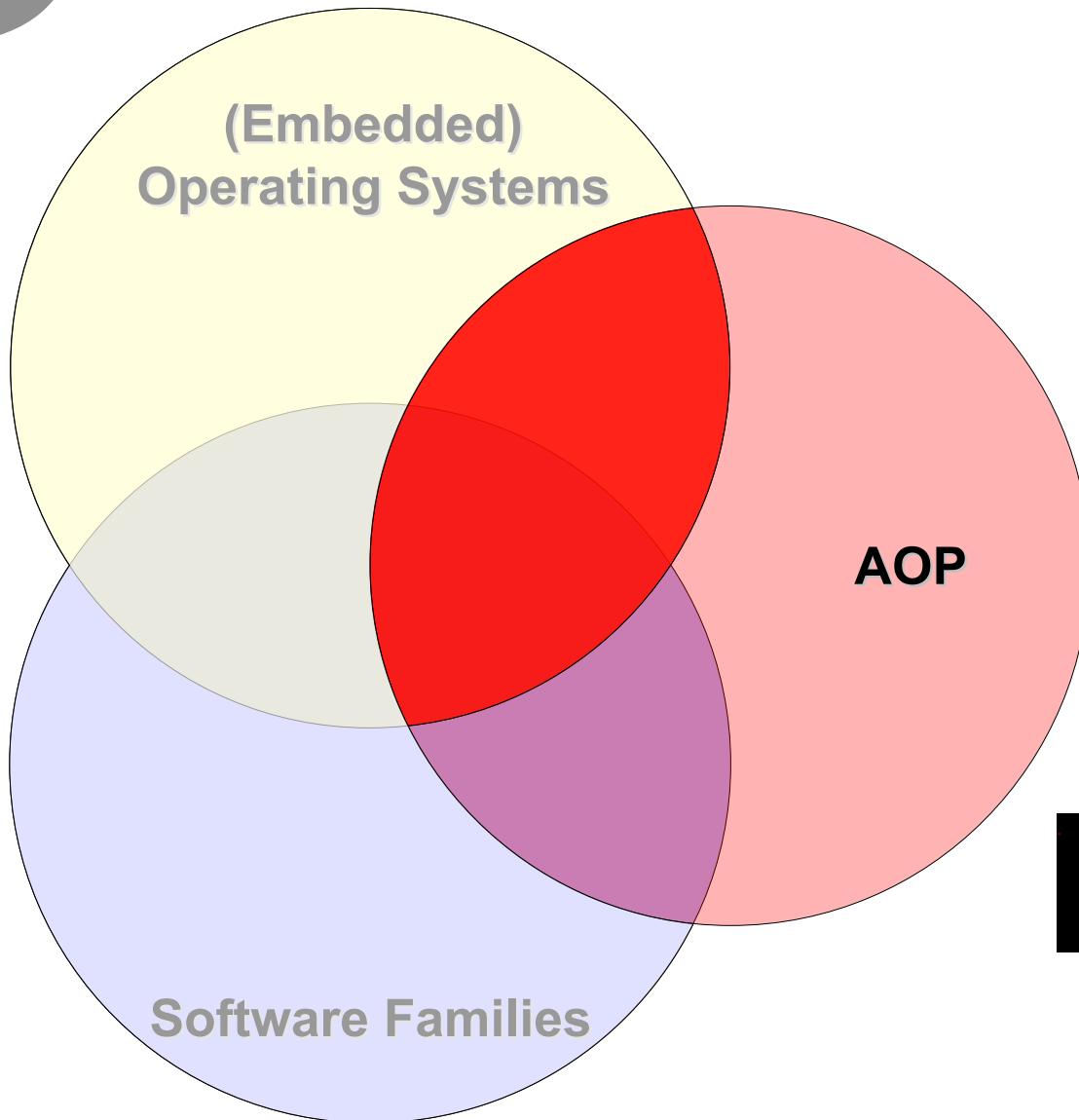
lohmann@informatik.uni-erlangen.de



Inhalt dieses Vortrags



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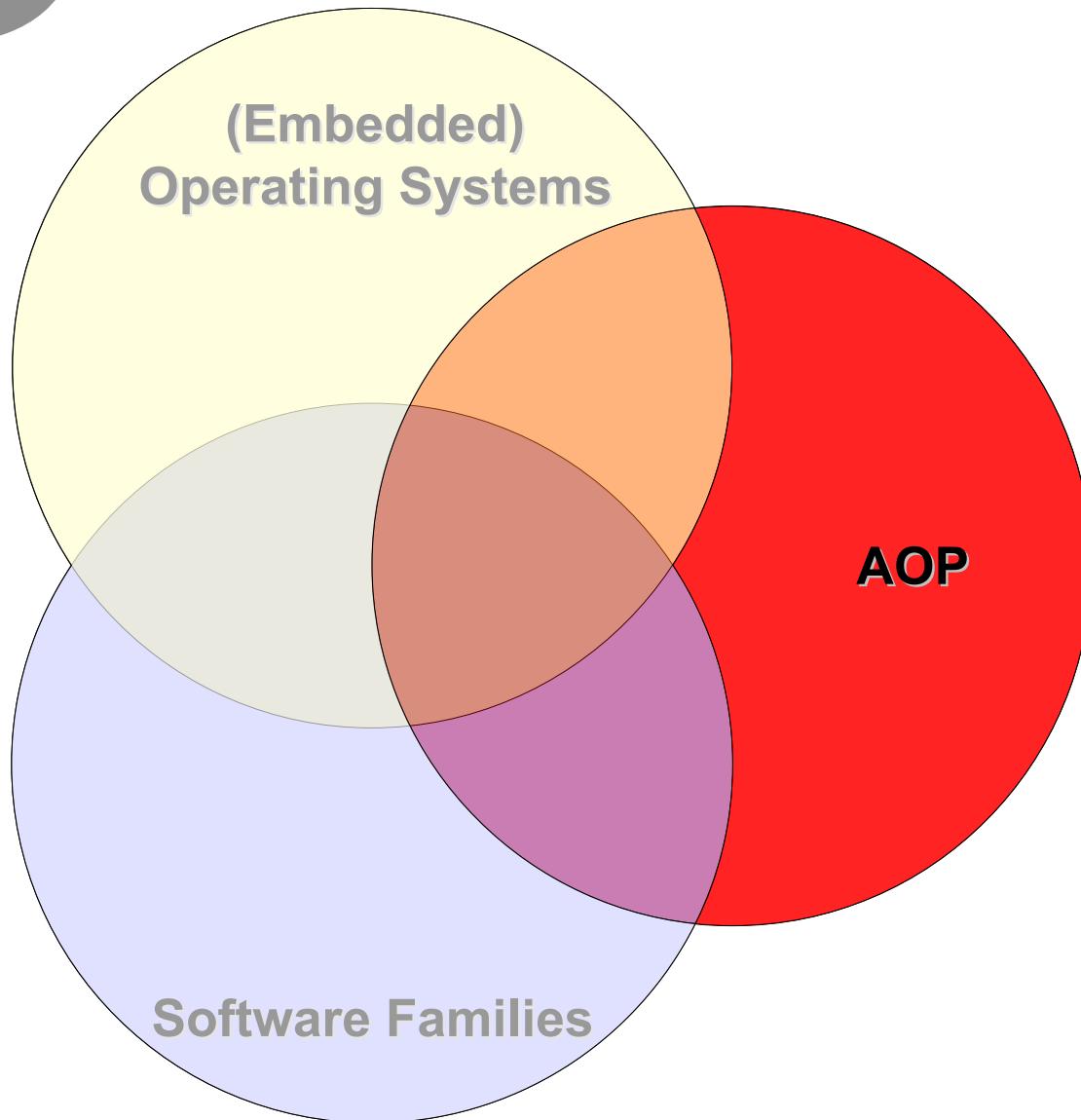


Motivation: Fallstudie eCos

- Das Problem
 - tangled code
 - scattered code
- AOP als Lösungsansatz



Inhalt dieses Vortrags



Hauptteil: AspectC++

- Die Sprache
 - Grundbegriffe
 - Besonderheiten
- Tutorial
 - "Queue" Familie



Überblick

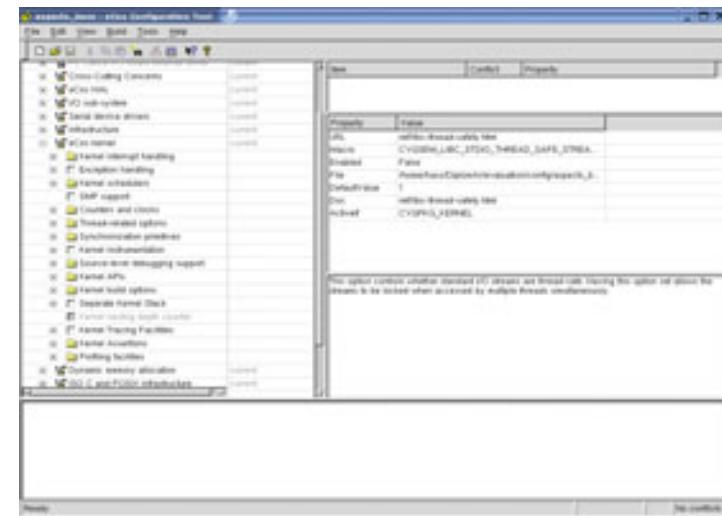
→ **Querschneidende Belange in eCos**

- Das Problem
- Aspektorientierte Programmierung
 - Der Lösungsansatz
- AspectC++
 - Grundlagen
 - Ergebnisse der eCos-Lösung
 - Tutorial
- Zusammenfassung

eCos: An OS Product Line



- Open source OS product line for embedded systems
 - developed and maintained by RedHat
 - supports a high number of 16/32 bit architectures
 - kernel written in C++
- Goal: static configurability
 - 63 selectable packages
 - 761 selectable configuration options
- Configuration approach
 - package selection (coarse-grained configuration)
 - **conditional compilation** (fine-grained configuration)

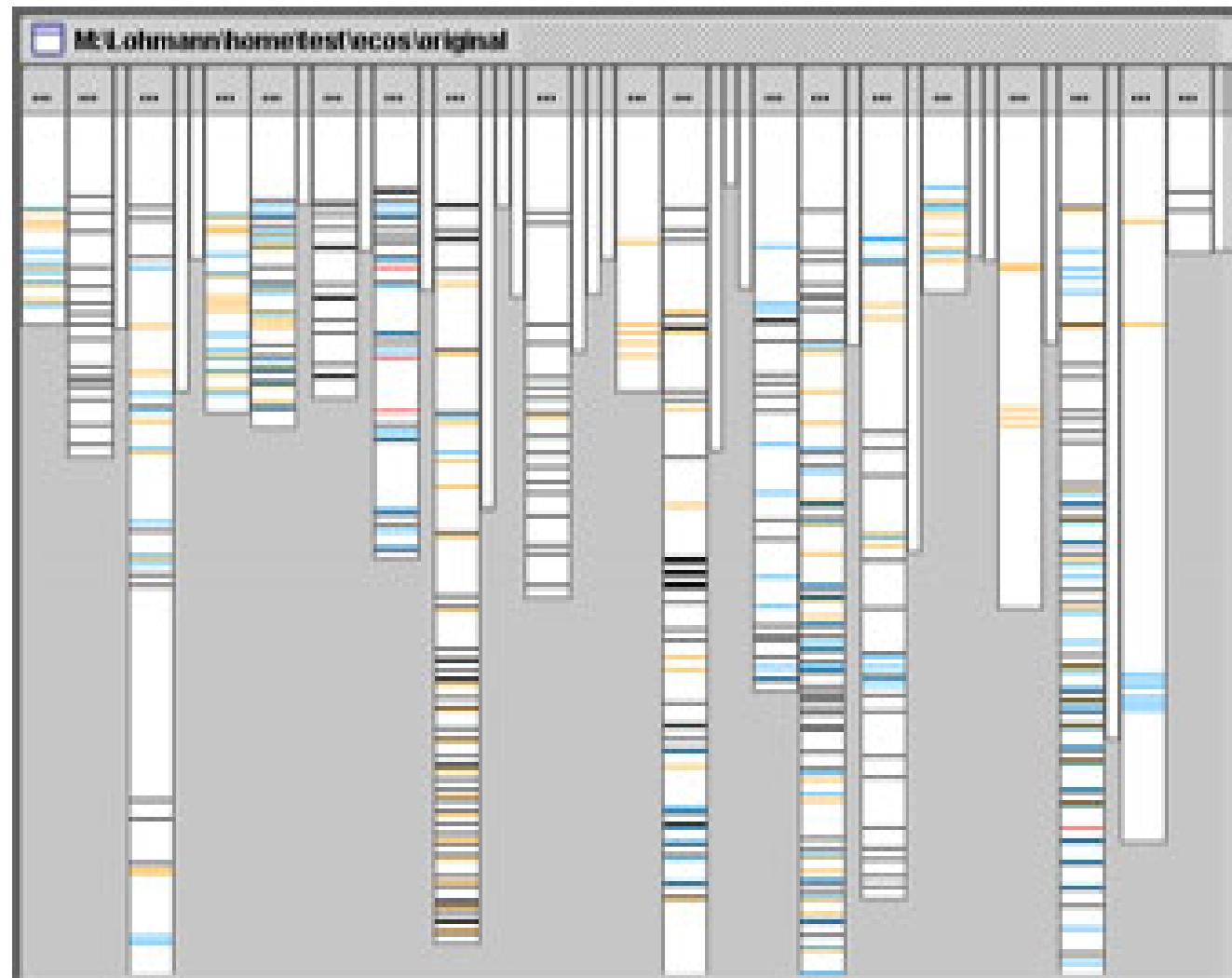


eCos: Global Policies in the Code

synchronization

instrumentation

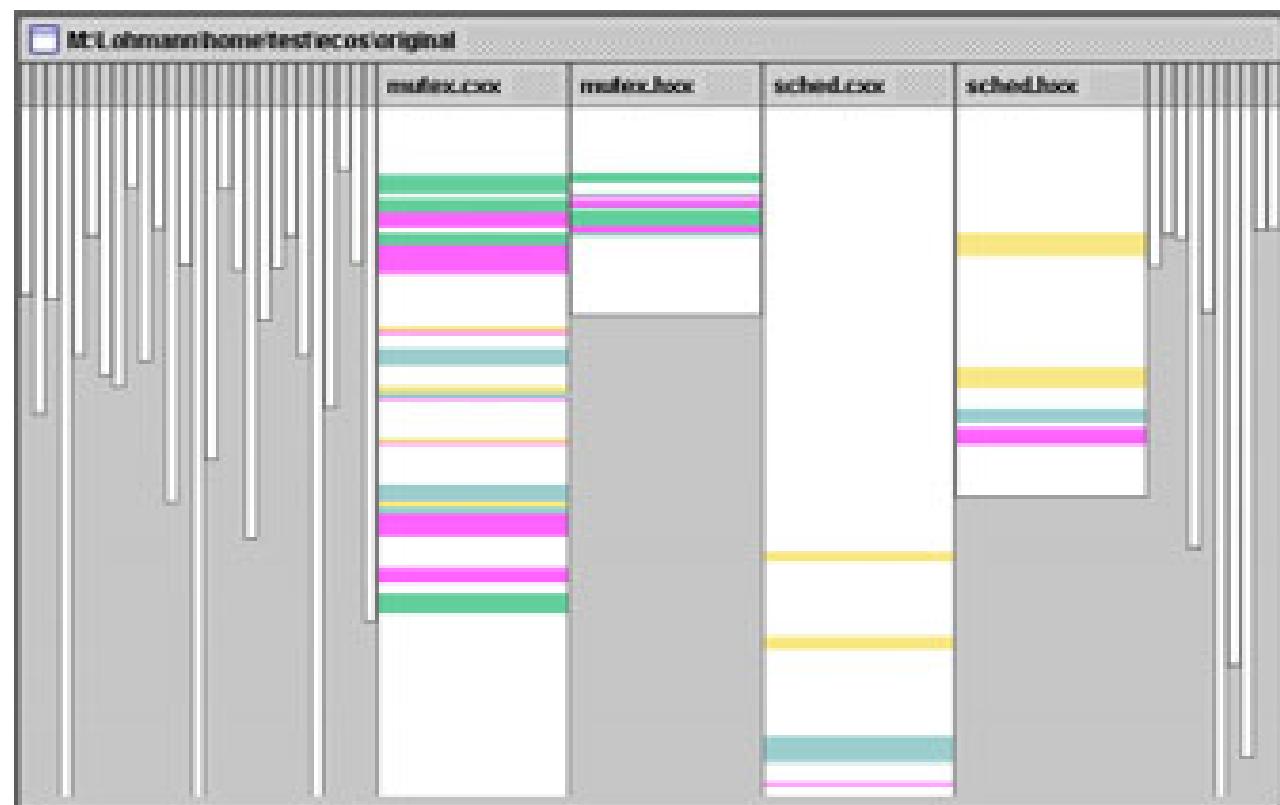
tracing



eCos: Configuration Options

Variants of the optional mutex priority inversion protocol

- simple
- ceiling
- inheritance
- dynamic



eCos: Implementation Example

```
Cyg_Mutex::Cyg_Mutex() {
    CYG_REPORT_FUNCTION();
    locked = false;
    owner = NULL;
#if defined(CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_DEFAULT) && \
    defined(CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_DYNAMIC)
#define CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_DEFAULT_INHERIT
    protocol = INHERIT;
#endif
#ifndef CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_DEFAULT_CEILING
    protocol = CEILING;
    ceiling = CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_DEFAULT_PRI;
#endif
#ifndef CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_DEFAULT_NONE
    protocol = NONE;
#endif
#else // not (DYNAMIC and DEFAULT defined)
#ifndef CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_CEILING
#ifndef CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_DEFAULT_PRIORITY
    // if there is a default priority ceiling defined, use that to initialize
    // the ceiling.
    ceiling = CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_DEFAULT_PRIORITY;
#else
    ceiling = 0; // Otherwise set it to zero.
#endif
#endif
#endif // DYNAMIC and DEFAULT defined
    CYG_REPORT_RETURN();
}
```

27 lines of code

eCos: Implementation Example

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    CYG_REPORT_RETURN();
}
```

2 lines for the tracing policy

eCos: Implementation Example

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#else
    ceiling = 0; // Otherwise set it to zero.
#endif
#endif
#endif // DYNAMIC and DEFAULT defined
    CYG_REPORT_RETURN();
}
```

21 (almost unreadable) lines for optional features

eCos: Implementation Example

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#else
    ceiling = 0; // Otherwise set it to zero.
#endif
#endif
#endif // DYNAMIC and DEFAULT defined
    CYG_REPORT_RETURN();
}
```

4 lines for the
basic implementation

eCos: Implementation Example

```
Cyg_Mutex::Cyg_Mutex() {
    CYG_REPORT_FUNCTION();
    locked = false;
    owner = NULL;
#if defined(CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_DEFAULT) && \
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    protocol = INHERIT;
#endif
#ifndef CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_DYNAMIC
    protocol = CEIL;
    ceiling = CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_DEFAULT;
#endif
#ifndef CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_NONE
    protocol = NONE;
#endif
#else // not (DYNAMIC or DEFAULT)
#define CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_DYNAMIC
#define CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_NONE
// if there is a ceiling, then it's dynamic
// the ceiling.
    ceiling = CYGSEM_KERNEL_SYNCH_MUTEX_PRIORITY_INVERSION_PROTOCOL_DYNAMIC;
#else
    ceiling = 0; // default
#endif
#endif // DYNAMIC and DEFAULT defined
    CYG_REPORT_RETURN();
}
```

Well, ...

- comprehensability ?
- extensability ?
- reuseability ?
- maintainability ?

Überblick

- Querschneidende Belange in eCos

- Das Problem



Aspektorientierte Programmierung

- Der Lösungsansatz

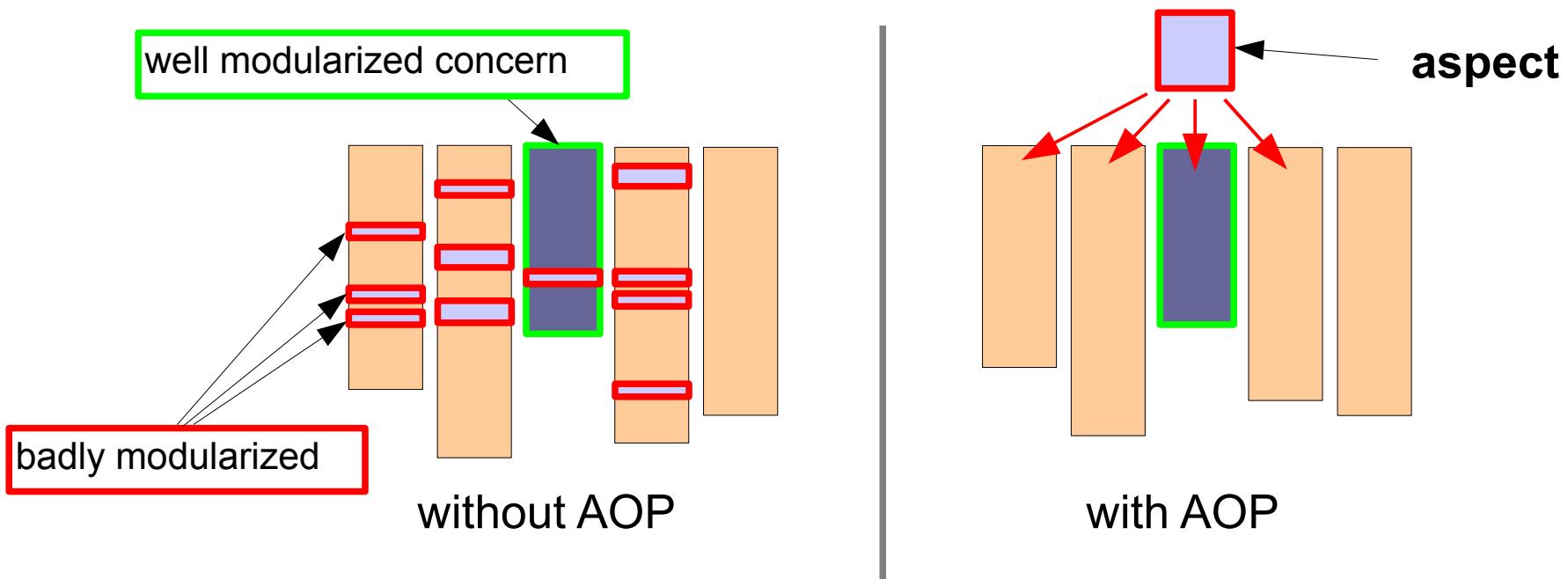
- AspectC++

- Grundlagen
 - Ergebnisse der eCos-Lösung
 - Tutorial

- Zusammenfassung

Aspekt-Orientierte Programmierung

AOP bietet Sprachmittel zur Trennung und Kapselung von querschneidenden Belangen



AOP: Trennung von *Was* und *Wo*

➤ Join-Points

- bezeichnen Stellen in der **Programmstruktur** (name join-points) oder Ereignisse im laufenden **Kontrollfluss** (code join-points)
- werden **deklarativ** spezifiziert durch **pointcut expressions**

➤ Advice

- zusätzliche **Elemente** (Methoden, Daten, ...), werden in bestimmte Klassen oder Strukturen eingefügt (statisch)
- zusätzliches **Verhalten** (Code), das
 - **before**, **after**
 - **around** (anstelle von)bestimmten Ereignissen im Kontrollfluss aktiviert wird

Überblick

- Querschneidende Belange in eCos
 - Das Problem
- Aspektorientierte Programmierung
 - Der Lösungsansatz

→ **AspectC++**

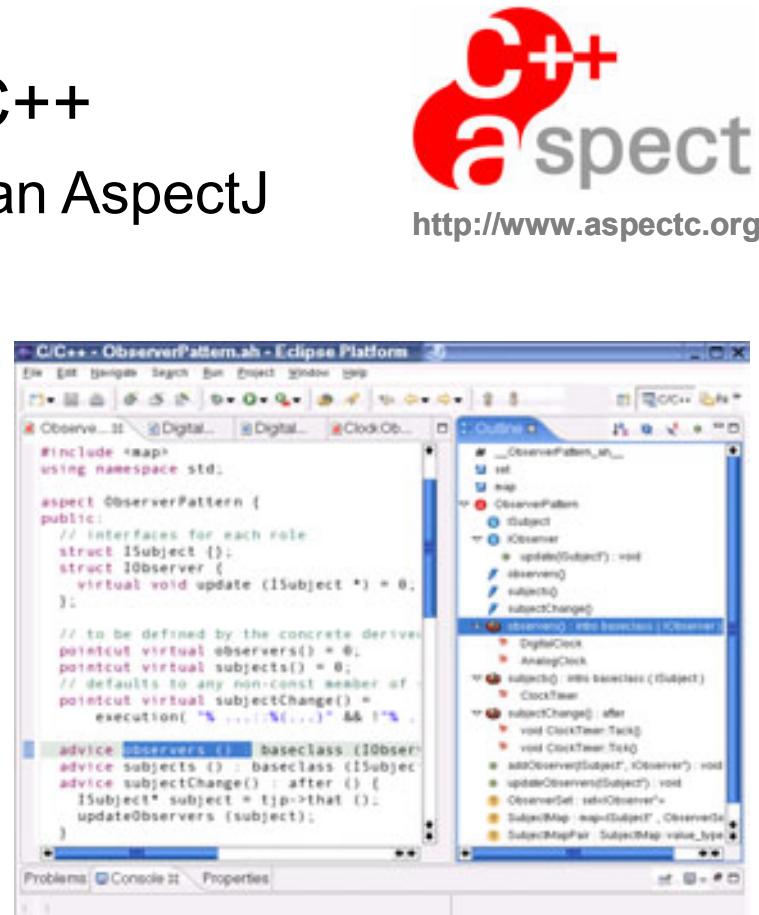
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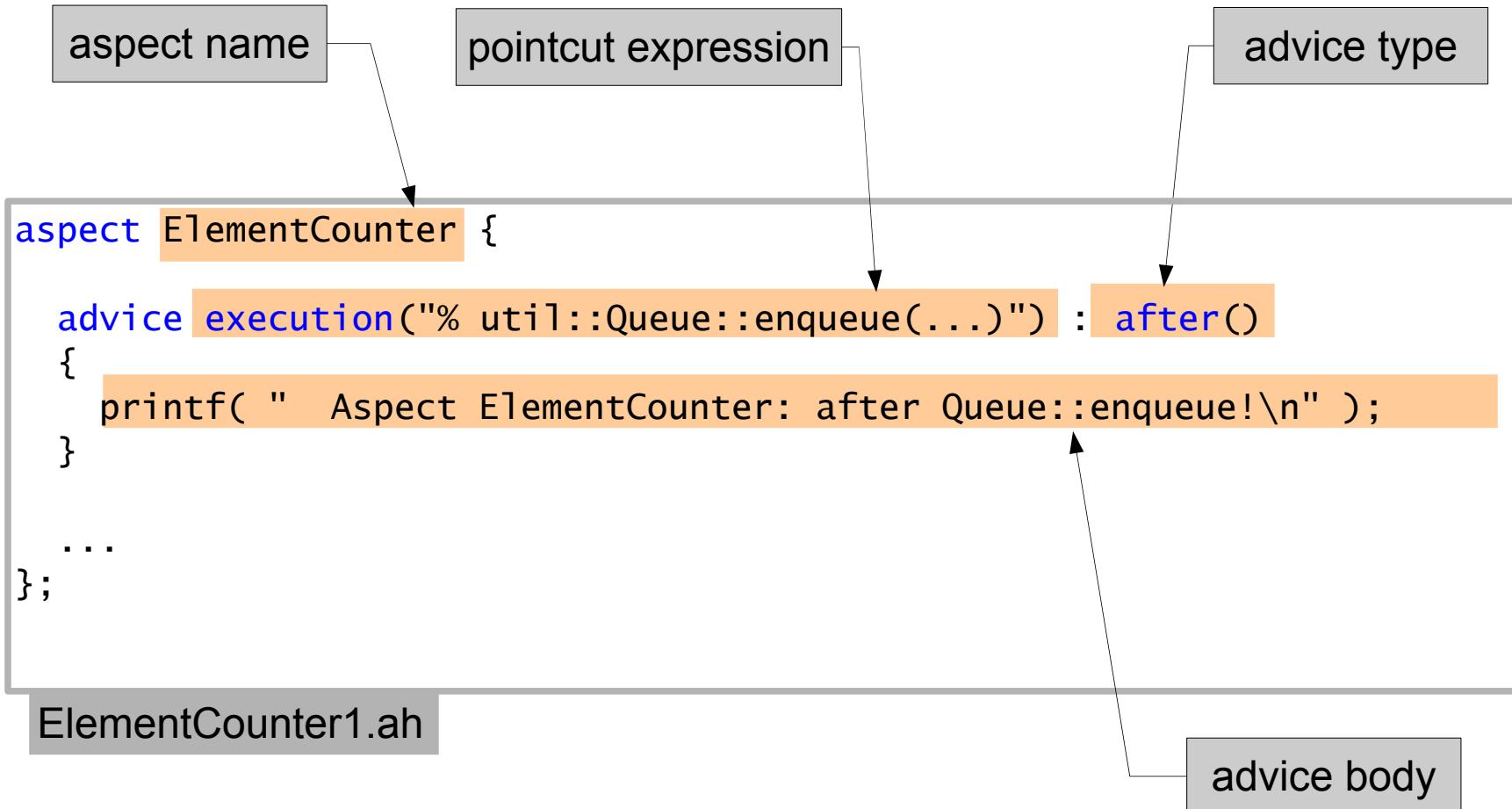
AspectC++

- AOP Spracherweiterung für C++
 - Syntax und Semantik angelehnt an AspectJ
- Werkzeugunterstützung
 - Source-Level Weber ac++
 - IDE Integration in Eclipse und (kommerziell) VisualStudio
 - OpenSource Projekt
 - Kommerzieller Support durch pure::systems GmbH
- Aktive Nutzer-Community





Syntactic Elements





Pointcut Expressions

- Pointcut expressions are made from ...
 - **match expressions**, e.g. "% util::queue::enqueue(...)"
 - are matched against C++ programm entities → name join-points
 - support wildcards
 - **pointcut functions**, e.g execution(...), call(...), that(...)
 - **execution**: all points in the control flow, where a function is about to be executed → code join-points
 - **call**: all points in the control flow, where a function is about to be called → code join-points
- Pointcut functions can be combined into expressions
 - using logical connectors: &&, ||, !
 - Example: `call("% util::Queue::enqueue(...)") && within("% main(...)")`



Advice

advice for code join-points (runtime events)

- **before advice**
 - advice code runs **before** the original code
 - may read and modify function parameter (call + execution)
- **after advice**
 - advice runs **after** the original code
 - may read and modify function result values (call + execution)
- **around advice**
 - advice code runs **instead of** the original code
 - original code can be explicitly invoked by `tjp->proceed()`

introductions

- further methods, attributes, ... are inserted into classes
- extension of interfaces and class implementations



Before / After Advice

for execution join points:

```
advice execution("void ClassA::foo()") : before()
```

```
advice execution("void ClassA::foo()") : after()
```

```
class ClassA {
public:
    void foo(){
        printf("ClassA::foo()\n");
    }
}
```

for call join points:

```
advice call ("void ClassA::foo()") : before()
```

```
advice call ("void ClassA::foo()") : after()
```

```
int main(){
    printf("main()\n");
    ClassA a;
    a.foo();
}
```



Around Advice

for execution join points:

```
advice execution("void ClassA::foo()") :  
around()  
  before code  
  
    tjp->proceed()  
  
  after code
```

```
class ClassA {  
public:  
  void foo(){  
    printf("ClassA::foo()\n");  
  }  
}
```

for call join points:

```
advice call("void ClassA::foo()") : around()  
  before code  
  
  tjp->proceed()  
  
  after code
```

```
int main(){  
  printf("main()\n");  
  ClassA a;  
  a.foo();  
}
```



Introductions

```
advice "ClassA" : slice class {  
    private element to introduce  
public:  
    public element to introduce  
};
```

```
class ClassA {  
public:  
    void foo(){  
        printf("ClassA::foo()\n");  
    }  
}
```

Beispiel: eCos Priority Ceiling Protocol

```
aspect priority_ceiling {  
  
    void call_clear_ceiling(Cyg_Thread*);  
    ...  
  
    advice "Cyg_Mutex" : cyg_priority_ceiling;  
    ...  
  
    advice construction("Cyg_Mutex") : after() {  
        tjp->that()->ceiling = CYGSEM_DEFAULT_PRIORITY;  
    }  
  
    advice call("% Cyg_Mutex::lock_inner(...)")  
        && within("% Cyg_Mutex::lock(...)")  
        && args(self)  
        : after(Cyg_Thread* self)  
    {  
        if(!(*tjp->result())) {  
            call_clear_ceiling(self);  
        }  
    }  
    ...  
};
```

Was Wo

Beispiel: eCos Priority Ceiling Protocol

```
aspect priority_ceiling {  
  
    void call_clear_ceiling(Cyg_Thread*);  
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    advice "Cyg_Mutex" : cyg_priority_ceiling;  
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        && within("% Cyg_Mutex::lock(...)")  
        && args(self)  
        : after(Cyg_Thread* self)  
    {  
        if(!(*tjp->result())) {  
            call_clear_ceiling(self);  
        }  
    }  
    ...  
};
```

Was

Wo

Einfügung einer Variable
ceiling in alle Klassen mit
dem Namen *Cyg_Mutex*

Beispiel: eCos Priority Ceiling Protocol

```

aspect priority_ceiling {

    void call_clear_ceiling(Cyg_Thread*);

    ...

    advice “Cyg_Mutex” : cyg_priority_ceiling;
    ...

    advice construction(“Cyg_Mutex”) : after() {
        tjp->that()->ceiling = CYGSEM_DEFAULT_PRIORITY;
    }

    advice call("% Cyg_Mutex::lock_inner(...)")
        && within("% Cyg_Mutex::lock(...)")
        && args(self)
        : after(Cyg_Thread* self)
    {
        if(!(*tjp->result())) {
            call_clear_ceiling(self);
        }
    }
    ...
};

    
```

Was

Wo

Ausführen der *Initialisierung*
nach dem **Erstellen** einer
Cyg_Mutex Instanz

Beispiel: eCos Priority Ceiling Protocol

```

aspect priority_ceiling {

    void call_clear_ceiling(Cyg_Thread*);

    ...

    advice “Cyg_Mutex” : cyg_priority_ceiling;
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    advice call("% Cyg_Mutex::lock_inner(...)")
        && within("% Cyg_Mutex::lock(...)")
        && args(self)
        : after(Cyg_Thread* self)
    {
        if(!(*tjp->result())) {
            call_clear_ceiling(self);
        }
    }
    ...
};

Was      Wo
};
```

Nach einem Aufruf von *Cyg_Mutex::lock_inner*, der stattfindet...

Beispiel: eCos Priority Ceiling Protocol

```

aspect priority_ceiling {

    void call_clear_ceiling(Cyg_Thread*);

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        && args(self)  

        : after(Cyg_Thread* self)
    {  

        if(!(*tjp->result())) {  

            call_clear_ceiling(self);
        }
    }
    ...
};

    
```

Was

Wo

Nach einem Aufruf von *Cyg_Mutex::lock_inner*, der stattfindet...

...innerhalb der Ausführung von *Cyg_Mutex::lock* und außerdem...

Beispiel: eCos Priority Ceiling Protocol

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aspect priority_ceiling {

    void call_clear_ceiling(Cyg_Thread*);

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    ...
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Was      Wo
};
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Nach einem Aufruf von *Cyg_Mutex::lock_inner*, der stattfindet...

...innerhalb der Ausführung von *Cyg_Mutex::lock* und außerdem...

...ein Argument vom Typ *Cyg_Thread** übergibt...

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Was           Wo

```

Nach einem Aufruf von *Cyg_Mutex::lock_inner*, der stattfindet...

...innerhalb der Ausführung von *Cyg_Mutex::lock* und außerdem...

...ein Argument vom Typ *Cyg_Thread** übergibt...

...überprüfe, ob der Mutex verlassen wurde und passe ggfs. die Priorität an

Example: Synchronization in AspeCos

ecos

```
aspect int_sync {  
  
    pointcut sync() = execution(...); // kernel calls to sync  
    || construction(...);  
    || destruction(...);  
  
    // advise kernel code to invoke lock() and unlock()  
    advice sync() : before() {  
        Cyg_Scheduler::lock();  
    }  
    advice sync() : after() {  
        Cyg_Scheduler::unlock();  
    }  
  
    // In eCos, a new thread always starts with a lock value of 0  
    advice execution(  
        "%Cyg_HardwareThread::thread_entry(...)") : before() {  
        Cyg_Scheduler::zero_sched_lock();  
    }  
    ...  
};
```



Example: Synchronization in AspeCos

ecos

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aspect int_sync {  
  
    pointcut sync() = execution(...); // kernel calls to sync  
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    }  
    ...  
};
```

where



Example: Synchronization in AspeCos

ecos

```
aspect int_sync {  
  
    pointcut sync() = execution(...); // kernel calls to sync  
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    || destruction(...);
```

where

```
// advise kernel code to invoke lock() and unlock()  
advice sync() : before() {  
    Cyg_Scheduler::lock();  
}  
advice sync() : after() {  
    Cyg_Scheduler::unlock();  
}
```

what

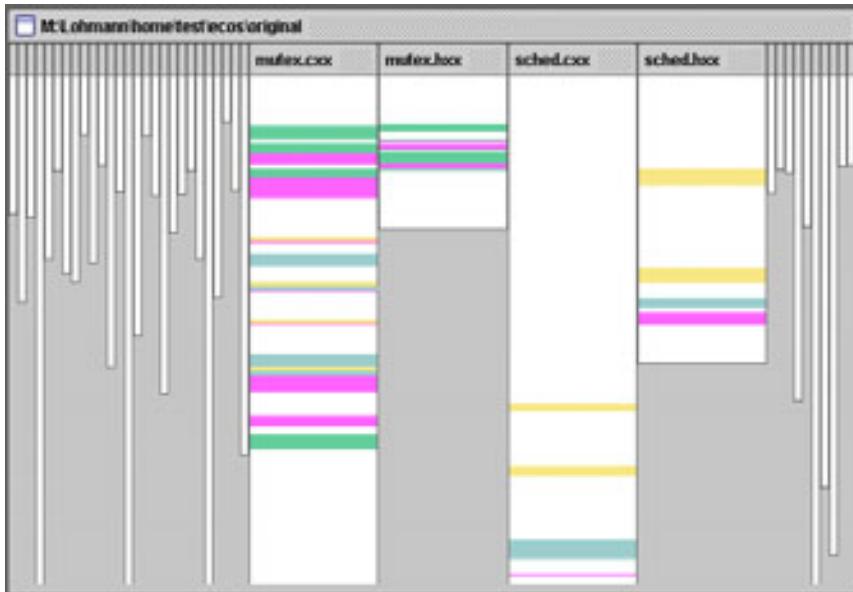
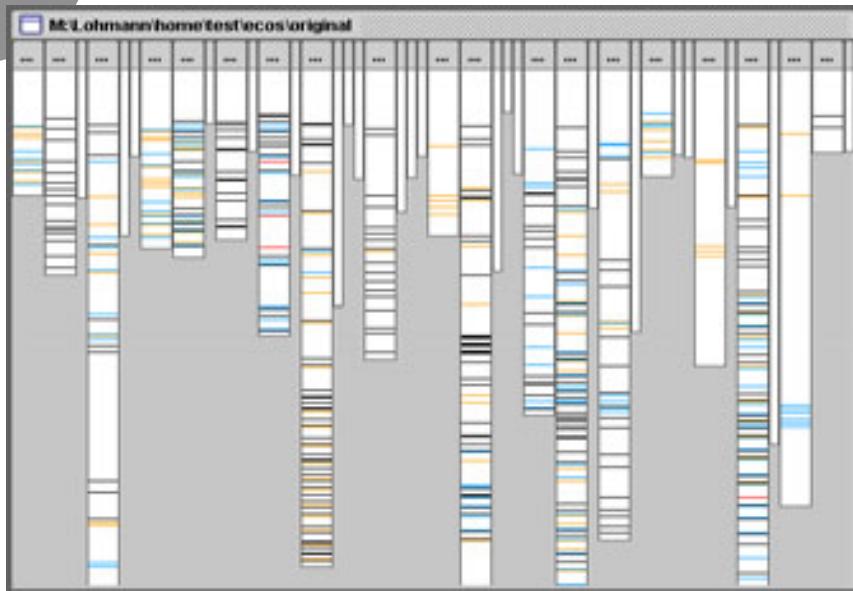
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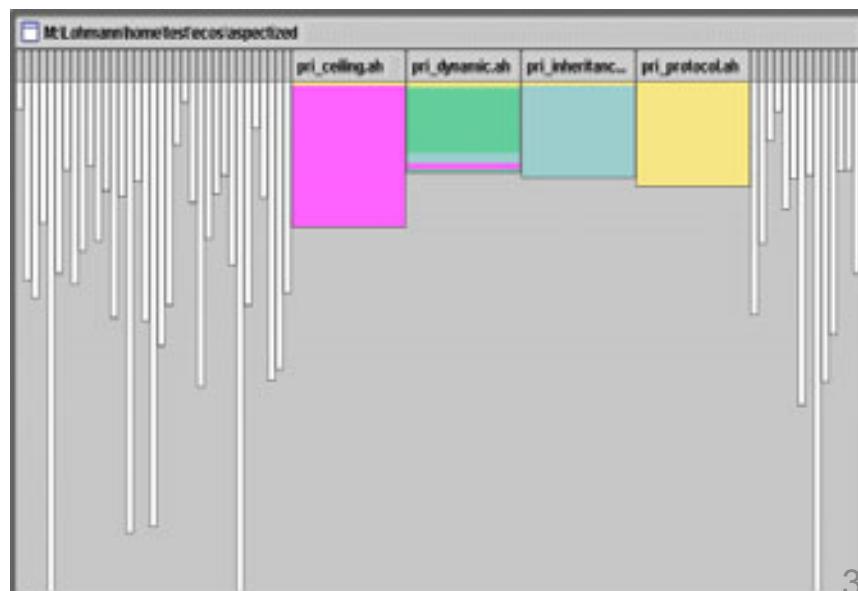
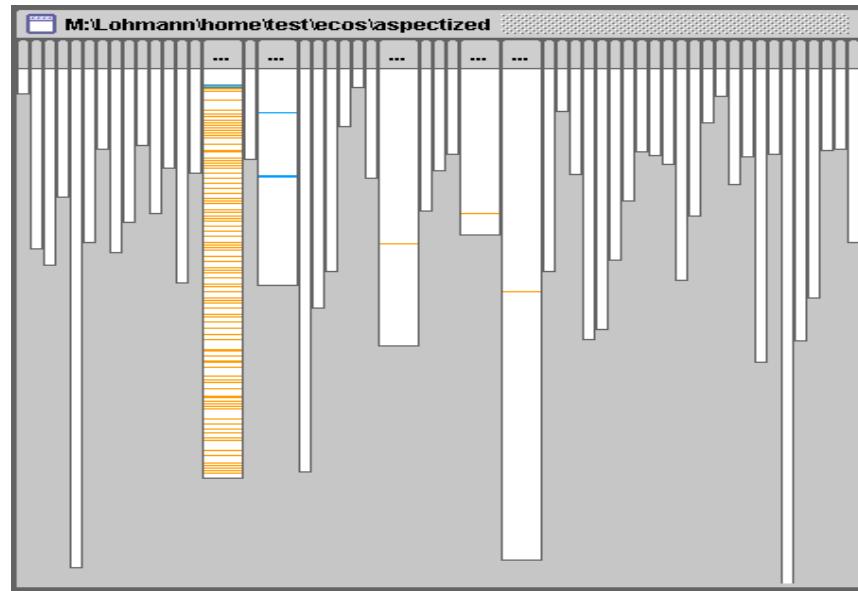
Results with eCos

- **Refactored:** *original kernel* → *aspectized kernel*
 - **3 cross-cutting policies**
 - interrupt synchronization 187 invocations → 160 code join-points
 - kernel instrumentation 162 invocations → 139 code join-points
 - tracing 336 invocations → 632 code join-points
 - **12 configuration options**
 - mutex features
 - thread features
- **Compared:** *original kernel* ↔ *aspectized kernel*
 - scattering, performance, memory footprint

original kernel



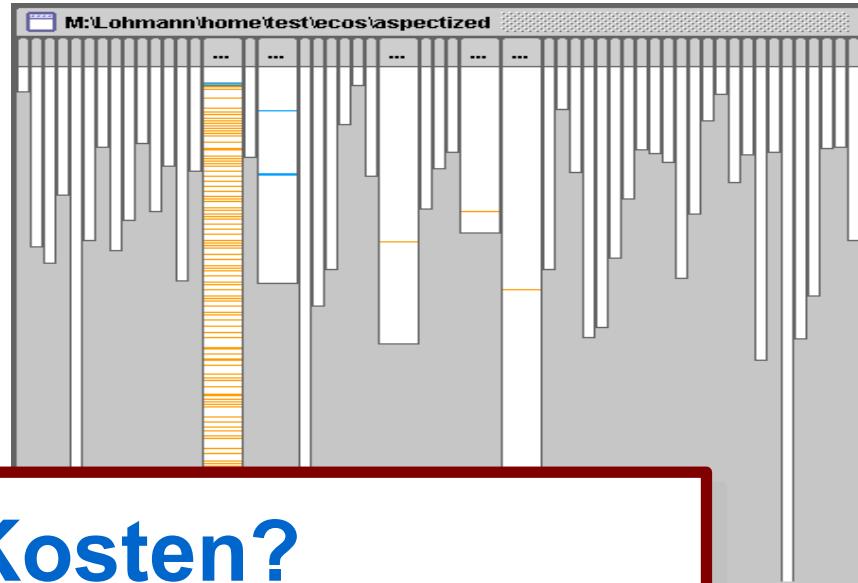
aspectized kernel



original kernel

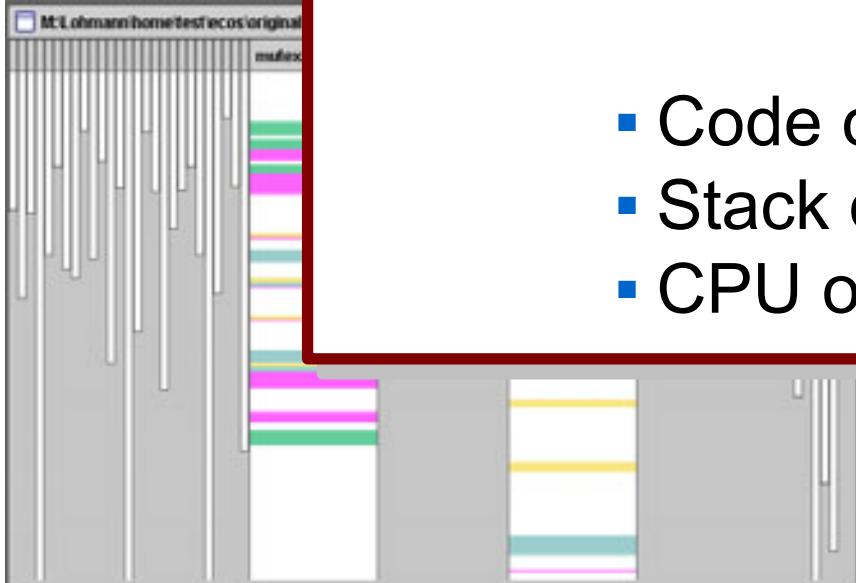


aspectized kernel



Und die Kosten?

- Code overhead: **0.9%**
- Stack overhead: **1.3%**
- CPU overhead: **0.0%**



Evaluation Case Study: CiAO-AS

CiAO



Specification of Operating System
V2.0.1

OS093: If interrupts are disabled and any OS services, excluding the interrupt services, are called outside of hook routines, then the Operating System shall return E_OS_DISABLEDINT



Evaluation Case Study: CiAO-AS

CiAO



Specification of Operating System
V2.0.1

OS093: If interrupts are disabled and any OS services, excluding the interrupt services, are called outside of hook routines, then the Operating System shall return E_OS_DISABLEDINT

```
aspect DisabledIntCheck {
    advice call( pc0SServices() && !pcInterruptServices() )
        && !within( pcHookRoutines() ) : around() {
        if( interruptsDisabled() )
            *tjp->result() = E_OS_DISABLEDINT;
        else
            tjp->proceed();
    } };
}
```



Evaluation Case Study: CiAO-AS

CiAO



Specification of Operating System
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aspect DisabledIntCheck {
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            *tjp->result() = E_OS_DISABLEDINT;
        else
            tjp->proceed();
    } };
}
```



Evaluation Case Study: CiAO-AS

CiAO



Specification of Operating System
V2.0.1

OS093: If interrupts are disabled and any OS services, **excluding the interrupt services**, are called outside of hook routines, then the Operating System shall return E_OS_DISABLEDINT

```
aspect DisabledIntCheck {
    advice call( pcOSServices() && !pcInterruptServices() )
        && !within( pcHookRoutines() ) : around() {
        if( interruptsDisabled() )
            *tjp->result() = E_OS_DISABLEDINT;
        else
            tjp->proceed();
    } };
}
```



Evaluation Case Study: CiAO-AS

CiAO



Specification of Operating System
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OS093: If interrupts are disabled and any OS services, excluding the interrupt services, are called outside of hook routines, then the Operating System shall return E_OS_DISABLEDINT

```
aspect DisabledIntCheck {
    advice call( pc0SServices() && !pcInterruptServices() )
        && !within( pcHookRoutines() ) : around() {
        if( interruptsDisabled() )
            *tjp->result() = E_OS_DISABLEDINT;
        else
            tjp->proceed();
    } };
}
```



Evaluation Case Study: CiAO-AS

CiAO



Specification of Operating System
V2.0.1

OS093: If interrupts are disabled and any OS services, excluding the interrupt services, are called outside of hook routines, then the Operating System shall return E_OS_DISABLEDINT

```
aspect DisabledIntCheck {
    advice call( pc0SServices() && !pcInterruptServices() )
        && !within( pcHookRoutines() ) : around() {
        if( interruptsDisabled() )
            *tjp->result() = E_OS_DISABLEDINT;
    else
        tjp->proceed();
    } };
}
```



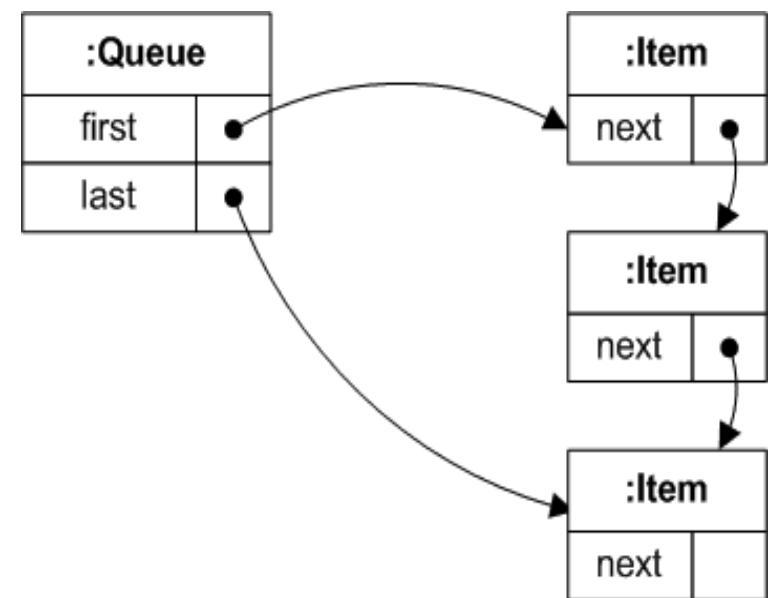
Überblick

- Querschneidende Belange in eCos
 - Das Problem
- Aspektorientierte Programmierung
 - Der Lösungsansatz
- AspectC++
 - Grundlagen
 - Ergebnisse der eCos-Lösung
 - **Tutorial**
- Zusammenfassung

Scenario: A Queue utility class

util::Queue
-first : util::Item
-last : util::Item
+enqueue(in item : util::Item)
+dequeue() : util::Item

util::Item
-next



The Simple Queue Class



```
namespace util {  
    class Item {  
        friend class Queue;  
        Item* next;  
    public:  
        Item() : next(0){}  
    };  
  
    class Queue {  
        Item* first;  
        Item* last;  
    public:  
        Queue() : first(0), last(0) {}  
  
        void enqueue( Item* item ) {  
            printf( " > Queue::enqueue()\n" );  
            if( last ) {  
                last->next = item;  
                last = item;  
            } else  
                last = first = item;  
            printf( " < Queue::enqueue()\n" );  
        }  
    };
```

```
    Item* dequeue() {  
        printf(" > Queue::dequeue()\n");  
        Item* res = first;  
        if( first == last )  
            first = last = 0;  
        else  
            first = first->next;  
        printf(" < Queue::dequeue()\n");  
        return res;  
    }  
}; // class Queue  
} // namespace util
```

Scenario: The Problem

Various users of Queue demand extensions:



I want Queue to throw exceptions!

Please extend the Queue class by an element counter!



Queue should be thread-safe!



The Not So Simple Queue Class



```
class Queue {  
    Item *first, *last;  
    int counter;  
    os::Mutex lock;  
public:  
    Queue () : first(0), last(0) {  
        counter = 0;  
    }  
    void enqueue(Item* item) {  
        lock.enter();  
        try {  
            if (item == 0)  
                throw QueueInvalidItemError();  
            if (last) {  
                last->next = item;  
                last = item;  
            } else { last = first = item; }  
            ++counter;  
        } catch (...) {  
            lock.leave(); throw;  
        }  
        lock.leave();  
    }
```

```
    Item* dequeue() {  
        Item* res;  
        lock.enter();  
        try {  
            res = first;  
            if (first == last)  
                first = last = 0;  
            else first = first->next;  
            if (counter > 0) -counter;  
            if (res == 0)  
                throw QueueEmptyError();  
        } catch (...) {  
            lock.leave();  
            throw;  
        }  
        lock.leave();  
        return res;  
    }  
    int count() { return counter; }  
}; // class Queue
```

What Code Does What?



```
class Queue {  
    Item *first, *last;  
    int counter;  
    os::Mutex lock;  
  
public:  
    Queue () : first(0), last(0) {  
        counter = 0;  
    }  
    void enqueue(Item* item) {  
        lock.enter();  
        try {  
            if (item == 0)  
                throw QueueInvalidItemError();  
            if (last) {  
                last->next = item;  
                last = item;  
            } else { last = first = item; }  
            ++counter;  
        } catch (...) {  
            lock.leave(); throw;  
        }  
        lock.leave();  
    }  
}
```

```
    Item* dequeue() {  
        Item* res;  
        lock.enter();  
        try {  
            res = first;  
            if (first == last)  
                first = last = 0;  
            else first = first->next;  
            if (counter > 0) -counter;  
            if (res == 0)  
                throw QueueEmptyError();  
        } catch (...) {  
            lock.leave();  
            throw;  
        }  
        lock.leave();  
        return res;  
    }  
    int count() { return counter; }  
}; // class Queue
```

The Simple Queue Class Revisited



```
namespace util {  
    class Item {  
        friend class Queue;  
        Item* next;  
    public:  
        Item() : next(0){}  
    };  
  
    class Queue {  
        Item* first;  
        Item* last;  
    public:  
        Queue() : first(0), last(0) {}  
  
        void enqueue( Item* item ) {  
            printf( "  > Queue::enqueue()\n" );  
            if( last ) {  
                last->next = item;  
                last = item;  
            } else  
                last = first = item;  
            printf( "  < Queue::enqueue()\n" );  
        }  
    };
```

```
    Item* dequeue() {  
        printf("  > Queue::dequeue()\n");  
        Item* res = first;  
        if( first == last )  
            first = last = 0;  
        else  
            first = first->next;  
        printf("  < Queue::dequeue()\n");  
        return res;  
    }  
}; // class Queue  
  
} // namespace util
```

Queue: Demanded Extensions



I. Element counting

Please extend
the Queue class
by an element
counter!

II. Errorhandling (signaling of errors by exceptions)

III. Thread safety (synchronization by mutex variables)



Element counting: The Idea

- Increment a counter variable after each execution of `util::Queue::enqueue()`
- Decrement it after each execution of `util::Queue::dequeue()`

ElementCounter1

```
aspect ElementCounter {  
  
    int counter;  
    ElementCounter() {  
        counter = 0;  
    }  
  
    advice execution("% util::Queue::enqueue(...)") : after() {  
        ++counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", counter );  
    }  
    advice execution("% util::Queue::dequeue(...)") : after() {  
        if( counter > 0 ) --counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", counter );  
    }  
};
```

ElementCounter1.ah

ElementCounter1 - Elements



```
aspect ElementCounter {  
    int counter;  
    ElementCounter() {  
        counter = 0;  
    }  
  
    advice execution("% util::Queue::enqueue(...)") : after() {  
        ++counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", counter );  
    }  
    advice execution("% util::Queue::dequeue(...)") : after() {  
        if( counter > 0 ) --counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", counter );  
    }  
};
```

We introduced a new **aspect** named *ElementCounter*.
An aspect starts with the keyword **aspect** and is syntactically much like a class.

ElementCounter1 - Elements



```
aspect ElementCounter {  
  
    int counter;  
    ElementCounter() {  
        counter = 0;  
    }  
  
    advice execution("% util::Queue::enqueue(...)") : after() {  
        ++counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", counter );  
    }  
    advice execution("% util::Queue::dequeue(...)") : after() {  
        if( counter > 0 ) --counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", counter );  
    }  
};
```

Like a class, an aspect can define data members, constructors and so on

ElementCounter1.ah

ElementCounter1 - Elements



```
aspect ElementCounter {  
    int counter;  
    ElementCounter() {  
        counter = 0;  
    }  
  
    advice execution("% util::Queue::enqueue(...)") : after() {  
        ++counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", counter );  
    }  
    advice execution("% util::Queue::dequeue(...)") : after() {  
        if( counter > 0 ) --counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", counter );  
    }  
};
```

We give **after advice** (= some crosscutting code to be executed after certain control flow positions)

A diagram consisting of two arrows originating from the 'advice' keyword in each of the two 'execution' blocks of the aspect code. These arrows point towards the explanatory text box on the right, which defines what 'after advice' means.

ElementCounter1.ah

ElementCounter1 - Elements

```
aspect ElementCounter {  
  
    int counter;  
    ElementCounter() {  
        counter = 0;  
    }  
  
    advice execution("% util::Queue::enqueue(...)") : after() {  
        ++counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", counter );  
    }  
    advice execution("% util::Queue::dequeue(...)") : after() {  
        if( counter > 0 ) --counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", counter );  
    }  
};
```

This **pointcut expression** denotes where the advice should be given.
(After **execution** of methods that match the pattern)

ElementCounter1.ah

ElementCounter1 - Elements



```
aspect ElementCounter {  
  
    int counter;  
    ElementCounter() {  
        counter = 0;  
    }  
  
    advice execution("% util::Queue::enqueue(...)") : after() {  
        ++counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", counter );  
    }  
    advice execution("% util::Queue::dequeue(...)") : after() {  
        if( counter > 0 ) --counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", counter );  
    }  
};
```

Aspect member elements can be accessed from within the advice body

ElementCounter1.ah

ElementCounter1 - Result



```
int main() {
    util::Queue queue;

    printf("main(): enqueueing an item\n");
    queue.enqueue( new util::Item );

    printf("main(): dequeuing two items\n");
    Util::Item* item;
    item = queue.dequeue();
    item = queue.dequeue();
}
```

main.cc

```
main(): enqueueing am item
> Queue::enqueue(00320FD0)
< Queue::enqueue(00320FD0)
Aspect ElementCounter: # of elements = 1
main(): dequeuing two items
> Queue::dequeue()
< Queue::dequeue() returning 00320FD0
Aspect ElementCounter: # of elements = 0
> Queue::dequeue()
< Queue::dequeue() returning 00000000
Aspect ElementCounter: # of elements = 0
<Output>
```

ElementCounter1 – What's next?



- The aspect is not the ideal place to store the counter, because it is shared between all Queue instances
- Ideally, counter becomes a member of Queue
- In the next step, we
 - move counter into Queue by **introduction**
 - **expose context** about the aspect invocation to access the current Queue instance

ElementCounter2

```
aspect ElementCounter {  
  
    advice "util::Queue" : slice class {  
        int counter;  
    public:  
        int count() const { return counter; }  
    };  
    advice execution("% util::Queue::enqueue(...)")  
        && that(queue) : after( util::Queue& queue ) {  
        ++queue.counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", queue.count() );  
    }  
    advice execution("% util::Queue::dequeue(...)")  
        && that(queue) : after( util::Queue& queue ) {  
        if( queue.count() > 0 ) --queue.counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", queue.count() );  
    }  
    advice construction("util::Queue")  
        && that(queue) : before( util::Queue& queue ) {  
        queue.counter = 0;  
    }  
};
```

ElementCounter2 - Elements



```
aspect ElementCounter {  
  
    advice "util::Queue" : slice class {  
        int counter;  
    public:  
        int count() const { return counter; }  
    };  
    advice execution("% util::Queue::enqueue(...)")  
        && that(queue) : after( util::Queue& queue ) {  
        ++queue.counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", queue.count() );  
    }  
    advice execution("% util::Queue::dequeue(...)")  
        && that(queue) : after( util::Queue& queue ) {  
        if( queue.count() > 0 ) --queue.counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", queue.count() );  
    }  
    advice construction("util::Queue")  
        && that(queue) : before( util::Queue& queue ) {  
        queue.counter = 0;  
    };  
};
```

Introduces a slice of members into all classes denoted by the pointcut "util::Queue"

ElementCounter2 - Elements



```
aspect ElementCounter {  
  
    advice "util::Queue" : slice class {  
        int counter;  
    public:  
        int count() const { return counter; }  
    };  
    advice execution("% util::Queue::enqueue(...)")  
        && that(queue) : after( util::Queue& queue ) {  
        ++queue.counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", queue.count() );  
    }  
    advice execution("% util::Queue::dequeue(...)")  
        && that(queue) : after( util::Queue& queue ) {  
        if( queue.count() > 0 ) --queue.counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", queue.count() );  
    }  
    advice construction("util::Queue")  
        && that(queue) : before( util::Queue& queue ) {  
        queue.counter = 0;  
    }  
};
```

We introduce a private *counter* element and a public method to read it

ElementCounter2 - Elements



```
aspect ElementCounter {  
    advice "util::Queue" : slice class {  
        int counter;  
    public:  
        int count() const { return counter; }  
    };  
    advice execution("% util::Queue::enqueue(...)")  
        && that(queue) : after(util::Queue& queue) {  
        ++queue.counter;  
        printf(" Aspect ElementCounter: # of elements = %d\n", queue.count());  
    }  
    advice execution("% util::Queue::dequeue(...)")  
        && that(queue) : after(util::Queue& queue) {  
        if( queue.count() > 0 ) --queue.counter;  
        printf(" Aspect ElementCounter: # of elements = %d\n", queue.count());  
    }  
    advice construction("util::Queue")  
        && that(queue) : before(util::Queue& queue) {  
        queue.counter = 0;  
    }  
};
```

A **context variable** `queue` is bound to *that* (the calling instance).
The calling instance has to be an `util::Queue`

ElementCounter2 - Elements



```
aspect ElementCounter {  
  
    advice "util::Queue" : slice class {  
        int counter;  
    public:  
        int count() const { return counter; }  
    };  
    advice execution("% util::Queue::enqueue(...)")  
        && that(queue) : after( util::Queue& queue ) {  
        ++queue.counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", queue.count() );  
    }  
    advice execution("% util::Queue::dequeue(...)")  
        && that(queue) : after( util::Queue& queue ) {  
        if( queue.count() > 0 ) --queue.counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", queue.count() );  
    }  
    advice construction("util::Queue")  
        && that(queue) : before( util::Queue& queue ) {  
        queue.counter = 0;  
    }  
};
```

The context variable `queue` is used to access the calling instance.

ElementCounter2 - Elements



```
aspect ElementCounter {  
  
    advice "util::Queue" : slice class {  
        int counter;  
    public:  
        int count() const { return counter; }  
    };  
    advice execution("% util::Queue::enqueue(...)")  
        && that(queue) : after( util::Queue& queue ) {  
        ++queue.counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", queue.count() );  
    }  
    advice execution("% util::Queue::dequeue(...)")  
        && that(queue) : after( util::Queue& queue ) {  
        if( queue.count() > 0 ) --queue.counter;  
        printf( " Aspect ElementCounter: # of elements = %d\n", queue.count() );  
    }  
    advice construction("util::Queue")  
        && that(queue) : before( util::Queue& queue ) {  
        queue.counter = 0;  
    }  
};
```

By giving **construction advice**
we ensure that counter gets
initialized

ElementCounter2 - Result



```
int main() {
    util::Queue queue;
    printf("main(): Queue contains %d items\n", queue.count());
    printf("main(): enqueueing some items\n");
    queue.enqueue(new util::Item);
    queue.enqueue(new util::Item);
    printf("main(): Queue contains %d items\n", queue.count());
    printf("main(): dequeuing one items\n");
    util::Item* item;
    item = queue.dequeue();
    printf("main(): Queue contains %d items\n", queue.count());
}
```

main.cc

ElementCounter2 - Result



```
int main() {
    util::Queue queue;
    printf("main(): Queue contains %d items\n", queue.count());
    printf("main(): enqueueing some items\n");
    queue.enqueue(new util::Item);
    queue.enqueue(new util::Item);
    printf("main(): Queue contains %d items\n", queue.count());
    printf("main(): dequeuing one item\n");
    util::Item* item;
    item = queue.dequeue();
    printf("main(): Queue contains %d items\n");
}
```

main.cc

```
main(): Queue contains 0 items
main(): enqueueing some items
> Queue::enqueue(00320FD0)
< Queue::enqueue(00320FD0)
Aspect ElementCounter: # of elements = 1
> Queue::enqueue(00321000)
< Queue]::enqueue(00321000)
Aspect ElementCounter: # of elements = 2
main(): Queue contains 2 items
main(): dequeuing one items
> Queue::dequeue()
< Queue::dequeue() returning 00320FD0
Aspect ElementCounter: # of elements = 1
main(): Queue contains 1 items
```

<Output>

Queue: Demanded Extensions



I. Element counting



I want Queue to
throw exceptions!

II. Errorhandling (signaling of errors by exceptions)

III. Thread safety (synchronization by mutex variables)

Errorhandling: The Idea

- We want to check the following constraints:
 - enqueue() is never called with a NULL item
 - dequeue() is never called on an empty queue
- In case of an error an exception should be thrown
- To implement this, we need access to ...
 - the parameter passed to enqueue()
 - the return value returned by dequeue()

... from within the advice

ErrorException

```
namespace util {
    struct QueueInvalidItemError {};
    struct QueueEmptyError {};
}

aspect ErrorException {

    advice execution("% util::Queue::enqueue(...)") && args(item)
        : before(util::Item* item) {
        if( item == 0 )
            throw util::QueueInvalidItemError();
    }

    advice execution("% util::Queue::dequeue(...)") && result(item)
        : after(util::Item* item) {
        if( item == 0 )
            throw util::QueueEmptyError();
    }
};
```

ErrorException.ah

ErrorException - Elements

```
namespace util {  
    struct QueueInvalidItemError {};  
    struct QueueEmptyError {};  
}  
  
aspect ErrorException {  
  
    advice execution("% util::Queue::enqueue(...)") && args(item)  
        : before(util::Item* item) {  
        if( item == 0 )  
            throw util::QueueInvalidItemError();  
    }  
    advice execution("% util::Queue::dequeue(...)") && result(item)  
        : after(util::Item* item) {  
        if( item == 0 )  
            throw util::QueueEmptyError();  
    }  
};
```

ErrorException.ah

We give advice to be executed *before* enqueue() and *after* dequeue()

ErrorException - Elements

```
namespace util {  
    struct QueueInvalidItemEr  
    struct QueueEmptyError {}  
}  
  
aspect ErrorException {  
  
    advice execution("% util::Queue::enqueue(...)") && args(item)  
        : before(util::Item* item) {  
        if( item == 0 )  
            throw util::QueueInvalidItemError();  
    }  
    advice execution("% util::Queue::dequeue(...)") && result(item)  
        : after(util::Item* item) {  
        if( item == 0 )  
            throw util::QueueEmptyError();  
    }  
};
```

ErrorException.ah

A context variable *item* is bound to the first argument of type *util::Item** passed to the matching methods

ErrorException - Elements

```
namespace util {  
    struct QueueInvalidItemEr  
    struct QueueEmptyError {}  
}  
  
aspect ErrorException {  
  
    advice execution("% util::Queue::enqueue(...)") && args(item)  
        : before(util::Item* item) {  
        if( item == 0 )  
            throw util::QueueInvalidItemError();  
    }  
    advice execution("% util::Queue::dequeue(...)") && result(item)  
        : after(util::Item* item) {  
        if( item == 0 )  
            throw util::QueueEmptyError();  
    }  
};
```

ErrorException.ah

Here the **context variable** *item* is bound to the **result** of type *util::Item** returned by the matching methods

Queue: Demanded Extensions



I. Element counting

Queue should be
thread-safe!

II. Errorhandling (signaling of errors by exceptions)



III. Thread safety (synchronization by mutex variables)

Thread Safety: The Idea



- Protect enqueue() and dequeue() by a mutex object
- To implement this, we need to
 - introduce a mutex variable into class Queue
 - lock the mutex before the execution of enqueue() / dequeue()
 - unlock the mutex after execution of enqueue() / dequeue()
- The aspect implementation should be exception safe!
 - in case of an exception, pending after advice is not called
 - solution: use around advice

LockingMutex

```
aspect LockingMutex {  
    advice "util::Queue" : slice class { os::Mutex lock; };  
  
    pointcut sync_methods() = "% util::Queue::%queue(...);  
  
    advice execution(sync_methods()) && that(queue)  
    : around( util::Queue& queue ) {  
        queue.lock.enter();  
        try {  
            tjp->proceed();  
        }  
        catch(...) {  
            queue.lock.leave();  
            throw;  
        }  
        queue.lock.leave();  
    }  
};
```

LockingMutex.ah

LockingMutex - Elements

```
aspect LockingMutex {  
    advice "util::Queue" : slice class { os::Mutex lock; };  
  
    pointcut sync_methods() = "% util::Queue::%queue(...);";  
  
    advice execution(sync_methods()) && that(queue)  
    : around( util::Queue& queue ) {  
        queue.lock.enter();  
        try {  
            tjp->proceed();  
        }  
        catch(...) {  
            queue.lock.leave();  
            throw;  
        }  
        queue.lock.leave();  
    }  
};
```

We introduce a mutex
member into class Queue

LockingMutex.ah

LockingMutex - Elements

```
aspect LockingMutex {  
    advice "util::Queue" : slice class { os::Mutex lock; };  
  
    pointcut sync_methods() = "% util::Queue::%queue(...);  
  
    advice execution(sync_methods()) && that(queue)  
    : around( util::Queue& queue ) {  
        queue.lock.enter();  
        try {  
            tjp->proceed();  
        }  
        catch(...) {  
            queue.lock.leave();  
            throw;  
        }  
        queue.lock.leave();  
    }  
};
```

Pointcuts can be named.
sync_methods describes all
methods that have to be
synchronized by the mutex

LockingMutex.ah

LockingMutex - Elements

```
aspect LockingMutex {  
    advice "util::Queue" : slice class { os::Mutex lock; };  
  
    pointcut sync_methods() = "% util::Queue::%queue(...);  
  
    advice execution(sync_methods()) && that(queue)  
    : around( util::Queue& queue ) {  
        queue.lock.enter();  
        try {  
            tjp->proceed();  
        }  
        catch(...) {  
            queue.lock.leave();  
            throw;  
        }  
        queue.lock.leave();  
    }  
};
```

sync_methods is used to give around advice to the execution of the methods

LockingMutex.ah

LockingMutex - Elements

```
aspect LockingMutex {  
    advice "util::Queue" : slice class { os::Mutex lock; };  
  
    pointcut sync_methods() = "% util::Queue::%queue(...);";  
  
    advice execution(sync_methods()) && that(queue)  
    : around( util::Queue& queue ) {  
        queue.lock.enter();  
        try {  
            tjp->proceed();  
        }  
        catch(...) {  
            queue.lock.leave();  
            throw;  
        }  
        queue.lock.leave();  
    }  
};
```

By calling `tjp->proceed()` the original method is executed

LockingMutex.ah

Queue: A new Requirement



- I. Element counting
- II. Errorhandling
(signaling of errors by exceptions)
- III. Thread safety
(synchronization by mutex variables)
- IV. Interrupt safety
(synchronization on interrupt level)

We need Queue to be synchronized on interrupt level!



Interrupt Safety: The Idea

- Scenario
 - Queue is used to transport objects between kernel code (interrupt handlers) and application code
 - If application code accesses the queue, interrupts must be disabled first
 - If kernel code accesses the queue, interrupts must not be disabled
- To implement this, we need to distinguish
 - if the call is made from kernel code, or
 - if the call is made from application code

LockingIRQ1

```
aspect LockingIRQ {  
  
    pointcut sync_methods() = "% util::Queue::%queue(...)";  
    pointcut kernel_code() = "% kernel::%(...)";  
  
    advice call(sync_methods()) && !within(kernel_code()) : around() {  
        os::disable_int();  
        try {  
            tjp->proceed();  
        }  
        catch(...) {  
            os::enable_int();  
            throw;  
        }  
        os::enable_int();  
    }  
};
```

LockingIRQ1.ah

LockingIRQ1 – Elements



```
aspect LockingIRQ {  
  
    pointcut sync_methods() = "% util::Queue::%queue(...);";  
    pointcut kernel_code() = "% kernel::%(...);"  
  
    advice call(sync_methods()) && !within(kernel_code()) : around() {  
        os::disable_int();  
        try {  
            tjp->proceed();  
        }  
        catch(...) {  
            os::enable_int();  
            throw;  
        }  
        os::enable_int();  
    }  
};
```

We define two pointcuts. One for the methods to be synchronized and one for all kernel functions

LockingIRQ1.ah

LockingIRQ1 – Elements



```
aspect LockingIRQ {  
  
    pointcut sync_methods() = "% util::Queue::%queue(...)";  
    pointcut kernel_code() = "% kernel::%(...)";  
  
    advice call(sync_methods()) && !within(kernel_code()) : around() {  
        os::disable_int();  
        try {  
            tjp->proceed();  
        }  
        catch(...) {  
            os::enable_int();  
            throw;  
        }  
        os::enable_int();  
    }  
};
```

This pointcut expression matches any call to a *sync_method* that is **not** done from *kernel_code*

LockingIRQ1.ah

Locking IRQ1 – Result

```

util::Queue queue;
void do_something() {
    printf("do_something()\n");
    queue.enqueue( new util::Item );
}
namespace kernel {
    void irq_handler() {
        printf("kernel::irq_handler()\n");
        queue.enqueue(new util::Item);
        do_something();
    }
}
int main() {
    printf("main()\n");
    queue.enqueue(new util::Item);
    kernel::irq_handler(); // irq
    printf("back in main()\n");
    queue.dequeue();
}

```

main.cc

```

main()
os::disable_int()
    > Queue::enqueue(00320FD0)
    < Queue::enqueue()
os::enable_int()
kernel::irq_handler()
    > Queue::enqueue(00321030)
    < Queue::enqueue()
do_something()
os::disable_int()
    > Queue::enqueue(00321060)
    < Queue::enqueue()
os::enable_int()
back in main()
os::disable_int()
    > Queue::dequeue()
    < Queue::dequeue() returning 00320FD0
os::enable_int()

```

<Output>

LockingIRQ1 – Problem

```

util::Queue queue;
void do_something() {
    printf("do_something()\n");
    queue.enqueue( new util::Item );
}
namespace kernel {
    void irq_handler() {
        printf("kernel::irq_handler()\n");
        queue.enqueue(new util::Item);
        do_something();
    }
}
int main() {
    printf("main()\n");
    queue.enqueue(new util::Item);
    kernel::irq_handler(); // irq
    printf("back in main()\n");
    queue.dequeue();
}

```

main.cc

The pointcut `within(kernel_code)`
does not match any **indirect** calls

ma
os to sync_methods

> Queue::enqueue(00320FD0)

< Queue::enqueue()

os::enable_int()

`kernel::irq_handler()`

> Queue::enqueue(00321030)

< Queue::enqueue()

`do_something()`

os::disable_int()

> Queue::enqueue(00321060)

< Queue::enqueue()

os::enable_int()

back in main()

os::disable_int()

> Queue::dequeue()

< Queue::dequeue() returning 00320FD0

os::enable_int()

<Output>

LockingIRQ2

```
aspect LockingIRQ {  
  
    pointcut sync_methods() = "% util::Queue::%queue(...);";  
    pointcut kernel_code() = "% kernel::%(...);";  
  
    advice execution(sync_methods())  
    && !cflow(execution(kernel_code())) : around() {  
        os::disable_int();  
        try {  
            tjp->proceed();  
        }  
        catch(...) {  
            os::enable_int();  
            throw;  
        }  
        os::enable_int();  
    }  
};
```

Solution

Using the **cflow** pointcut function

LockingIRQ2.ah

LockingIRQ2 – Elements

```

aspect LockingIRQ {

    pointcut sync_methods() = "% util::Queue::%queue(...)";
    pointcut kernel_code() = "% kernel::%(...);"

    advice execution(sync_methods())
    && !cflow(execution(kernel_code())) : around() {
        os::disable_int();
        try {
            tjp->proceed();
        }
        catch(...) {
            os::enable_int();
            throw;
        }
        os::enable_int();
    }
};

```

This pointcut expression matches the execution of *sync_methods* if no *kernel_code* is on the call stack. *cflow* checks the call stack (control flow) at runtime.

LockingIRQ2.ah

Locking IRQ2 – Result



```
util::Queue queue;
void do_something() {
    printf("do_something()\n");
    queue.enqueue( new util::Item );
}
namespace kernel {
    void irq_handler() {
        printf("kernel::irq_handler()\n");
        queue.enqueue(new util::Item);
        do_something();
    }
}
int main() {
    printf("main()\n");
    queue.enqueue(new util::Item);
    kernel::irq_handler(); // irq
    printf("back in main()\n");
    queue.dequeue();
}
```

main.cc

```
main()
os::disable_int()
    > Queue::enqueue(00320FD0)
    < Queue::enqueue()
os::enable_int()
kernel::irq_handler()
    > Queue::enqueue(00321030)
    < Queue::enqueue()
do_something()
    > Queue::enqueue(00321060)
    < Queue::enqueue()
back in main()
os::disable_int()
    > Queue::dequeue()
    < Queue::dequeue() returning 00320FD0
os::enable_int()
```

<Output>

Überblick

- **Querschneidende Belange in eCos**
 - Das Problem
- **Aspektorientierte Programmierung**
 - Der Lösungsansatz
- **AspectC++**
 - Grundlagen
 - Ergebnisse der eCos-Lösung
 - Tutorial
- **Zusammenfassung**



AspectC++: Advanced Concepts



- Join Point API
 - provides a uniform interface to the aspect invocation context, both at runtime and compile-time
- Abstract Aspects and Aspect Inheritance
 - comparable to class inheritance, aspect inheritance allows to reuse parts of an aspect and overwrite other parts
- Generic Advice
 - exploits static type information in advice code
- Aspect Ordering
 - allows to specify the invocation order of multiple aspects
- Aspect Instantiation
 - allows to implement user-defined aspect instantiation models

The Joinpoint API

- Inside an advice body, the current joinpoint context is available via the **implicitly passed tjp** variable:

```
advice ... {
    struct JoinPoint {
        ...
        } *tjp;      // implicitly available in advice code
        ...
}
```

- You have already seen how to use **tjp**, to ...
 - execute the original code in around advice with **tjp->proceed()**
- The joinpoint API provides a rich interface
 - to expose context **independently** of the aspect target
 - this is especially useful in writing **reusable aspect code**

The Join Point API (Excerpt)



Types (compile-time)

```
// object type (initiator)  
That  
  
// object type (receiver)  
Target  
  
// result type of the affected function  
Result  
  
// type of the i'th argument of the affected  
// function (with 0 <= i < ARGS)  
Arg<i>::Type  
Arg<i>::ReferredType
```

Consts (compile-time)

```
// number of arguments  
ARGS  
  
// unique numeric identifier for this join point  
JPID  
  
// numeric identifier for the type of this join  
// point (AC::CALL, AC::EXECUTION, ...)  
JPTYPE
```

Values (runtime)

```
// pointer to the object initiating a call  
That* that()  
  
// pointer to the object that is target of a call  
Target* target()  
  
// pointer to the result value  
Result* result()  
  
// typed pointer the i'th argument value of a  
// function call (compile-time index)  
Arg<i>::ReferredType* arg()  
  
// pointer the i'th argument value of a  
// function call (runtime index)  
void* arg( int i )  
  
// textual representation of the joinpoint  
// (function/class name, parameter types...)  
static const char* signature()  
  
// executes the original joinpoint code  
// in an around advice  
void proceed()  
  
// returns the runtime action object  
AC::Action& action()
```

Abstract Aspects and Inheritance



- Aspects can inherit from other aspects...
 - Reuse aspect definitions
 - Override methods and pointcuts
- Pointcuts can be pure virtual
 - Postpone the concrete definition to derived aspects
 - An aspect with a pure virtual pointcut is called **abstract aspect**
- Common usage: Reusable aspect implementations
 - Abstract aspect defines advice code, but pure virtual pointcuts
 - Aspect code uses the joinpoint API to expose context
 - Concrete aspect inherits the advice code and overrides pointcuts

Abstract Aspects and Inheritance



```
#include "mutex.h"
aspect LockingA {
    pointcut virtual sync_classes() = 0;
    pointcut virtual sync_methods() = 0;

    advice sync_classes() : slice class {
        os::Mutex lock;
    };
    advice execution(sync_methods()) : around() {
        tjp->that()->lock.enter();
        try {
            tjp->proceed();
        }
        catch(...) {
            tjp->that()->lock.leave();
            throw;
        }
        tjp->that()->lock.leave();
    };
};
```

LockingA.ah

The abstract locking aspect declares two **pure virtual pointcuts** and uses the **joinpoint API** for an context-independent advice implementation.

```
#include "LockingA.ah"

aspect LockingQueue : public LockingA {
    pointcut sync_classes() =
        "util::Queue";
    pointcut sync_methods() =
        "% util::Queue::%queue(...)";
};
```

LockingQueue.ah

Abstract Aspects and Inheritance



```
#include "mutex.h"
aspect LockingA {
    pointcut virtual sync_classes() = 0;
    pointcut virtual sync_methods() = 0;

    advice sync_classes() : slice class {
        os::Mutex lock;
    };
    advice execution(sync_methods()) : around() {
        tjp->that()->lock.enter();
        try {
            tjp->proceed();
        }
        catch(...) {
            tjp->that()->lock.leave();
            throw;
        }
        tjp->that()->lock.leave();
    }
};
```

LockingA.ah

The concrete locking aspect **derives** from the abstract aspect and **overrides** the pointcuts.

```
#include "LockingA.ah"

aspect LockingQueue : public LockingA {
    pointcut sync_classes() =
        "util::Queue";
    pointcut sync_methods() =
        "% util::Queue::%queue(...)";
};
```

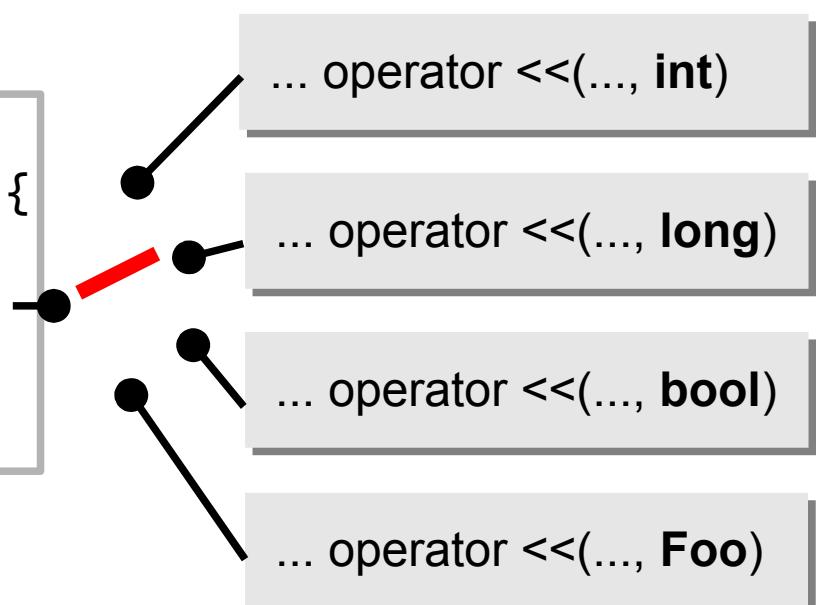
LockingQueue.ah

Generic Advice

Uses static JP-specific type information in advice code

- in combination with C++ overloading
- to instantiate C++ templates and template meta-programs

```
aspect TraceService {  
    advice call(...) : after() {  
        ...  
        cout << *tjp->result();  
    }  
};
```



Generic Advice

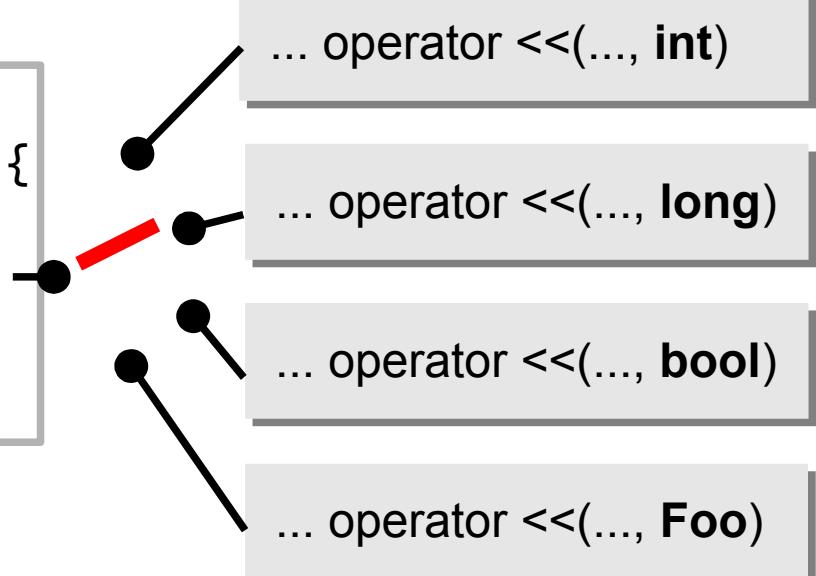
Uses static JP-specific type information in advice code

- in combination with C++ overloading

Resolves to the **statically typed** return value template meta-programs

- no runtime type checks are needed
- unhandled types are detected at compile-time
- functions can be inlined

```
aspect TraceService {  
    advice call(...) : after() {  
        ...  
        cout << *tjp->result();  
    }  
};
```



Aspect Ordering

- Aspects should be independent of other aspects
 - However, sometimes inter-aspect dependencies are unavoidable
 - Example: Locking should be activated before any other aspects
- Order advice
 - The aspect order can be defined by **order advice**
`advice pointcut-expr : order(high, ..., low)`
 - Different aspect orders can be defined for different pointcuts
- Example

```
advice "% util::Queue::%queue(...)"  
: order( "LockingIRQ", "%" && !"LockingIRQ" );
```

Aspect Instantiation



- Aspects are singletons by default
 - **aspectof()** returns pointer to the one-and-only aspect instance
- By overriding **aspectof()** this can be changed
 - e.g. one instance per client or one instance per thread

```
aspect MyAspect {
    // ...
    static MyAspect* aspectof() {
        static __declspec(thread) MyAspect* theAspect;
        if( theAspect == 0 )
            theAspect = new MyAspect;
        return theAspect;
    }
};
```

MyAspect.ah

Example of an user-defined aspectof() implementation for per-thread aspect instantiation by using thread-local storage.

(Visual C++)

Überblick

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 - Der Lösungsansatz
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 - Tutorial

 **Zusammenfassung**

Zusammenfassung

- AOP ...
 - versucht das Problem der querschneidenden Belange zu lösen: Vermeidung von *tangling* und *scattering*
 - modulare Implementierung durch Aspekte
 - Trennung von WO und WAS
- AspectC++ ...
 - erlaubt AOP mit C++
 - ähnelt AspectJ
 - wird durch IDEs unterstützt
- Beispiele zeigen beispielsweise, dass ...
 - konfigurierbare Software von Aspekten profitiert
 - redundanter Code vermieden werden kann

Weitere Informationen

- **das Web Portal der Community:** www-aosd.net
 - weitere AOP Sprachen/Werkzeuge
 - Konferenzen/Workshops
 - mailing lists
- **AspectC++:** www.aspectc.org
 - alle Infos zum AspectC++ Projekt
 - mailing list
- **Literatur:** "Aspect-Oriented Software Development"
 - von R. Filman, T. Elrad, S. Clarke, M. Aksit



Vielen Dank für's zuhören!